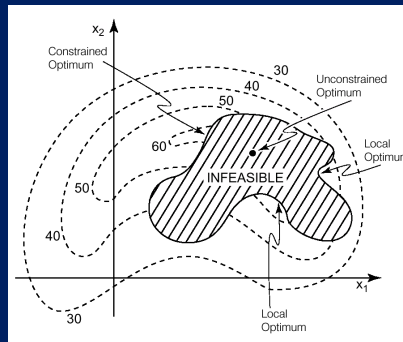


# A New Method For Numerical Constrained Optimization



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## Motivation

- The applicability of optimization methods is widespread, reaching into almost *every* activity in which numerical information is processed
- For a summary of applications and theory
  - ☞ See Fletcher “Practical Methods of Optimization”
- For numerous applications in computer graphics
  - ☞ See Goldsmith and Barr “Applying constrained optimization to computer graphics”
- In this sketch, we describe a method and not its application

## Informal Problem Statement

- An ideal problem for constrained optimization
  - ☞ has a single measure defining the quality of a solution (called the *objective function*  $F$ )
  - ☞ plus some requirements upon that solution that must not be violated (called the *constraints*  $C_i$ )
- A constrained optimization method maximizes (or minimizes)  $F$  while satisfying the  $C_i$ 's
- Both  $F$  and  $C_i$ 's are functions of  $\mathbf{x} \in \mathbb{R}^N$ , the input parameters to be determined



## Informal Problem Statement

- Many flavors of optimization
  - ☞  $\mathbf{x}$  can be real-valued, integer, mixed
  - ☞  $F$  and  $C_i$ 's can be linear, quadratic, nonlinear
  - ☞  $F$  and  $C_i$ 's can be smooth (i.e., differentiable) or nonsmooth
  - ☞  $F$  and  $C_i$ 's can be noisy or noise-free
  - ☞ methods can be globally convergent or global
- Our focus
  - ☞ globally convergent methods
  - ☞ real-valued, nonlinear, potentially nonsmooth, potentially noisy, constrained problems



## Our Contribution

- A new method for constraint handling, called *partitioned performances*, that
  - ☞ can be applied to established optimization algorithms
  - ☞ can improve their ability to traverse constrained space
- A new optimization method, called *SPIDER*, that
  - ☞ applies partitioned performances to a new variation of the Nelder and Mead polytope algorithm



## An observation leads to an idea

- Observation
  - ☞ Many constrained problems have optima that lie near constraint boundaries
  - ☞ Consequently, avoidance (or approximations) of constraints can hinder an algorithm's path to the answer
- Idea
  - ☞ By allowing (and even *encouraging*) an optimization algorithm to move its vertices into constrained space, a more efficient and robust algorithm emerges



## The idea leads to a method

- Constraints are partitioned (i.e., grouped) into multiple levels (i.e., categories)
- A constrained performance, independent of the objective function, is defined for each level
- A set of rules, based on these *partitioned performances*, specify the ordering and movement of vertices as they straddle constraint boundaries
- These rules are non-greedy, permitting vertices at a higher (i.e., better) level to move to a lower (i.e., worse) level



## Partitioned Performances (Advantages)

- Do not use a penalty function and thus do not warp the performance surface
  - ☞ this avoids the possible ill-conditioning of the objective function typical in penalty methods
- Do not linearize the constraints as do other methods (e.g., SQP)
- Assume very little about the problem form
  - ☞  $F$  and  $C_i$ 's can be nonsmooth (i.e., nondifferentiable) and highly nonlinear



## Partitioning Constraints

- One effective partitioning of constraints
  - ☞ place simple limits on  $\mathbf{x} \in \mathbb{R}^N$  into level 1 (e.g.,  $x_1 \geq 0$ )
  - ☞ place constraints which, when violated, produce singularities in  $F$  into level 1
  - ☞ all other constraints into level 2
  - ☞ and the objective function  $F$  into level 3
- Many different strategies for partitioning
  - ☞ just two levels: constrained and feasible
  - ☞ a level for every constraint, and a feasible level
  - ☞ *dynamic* partitioning (changing the level assignments during the search)



## Computing Performance

- Assume a partitioning of  $F$  and the  $C_i$ 's into  $W$  levels  $[L_1 \dots L_w]$  with  $L_w = \{ F \}$
- We define the *partitioned performance* of a location  $\mathbf{x} \in \mathbb{R}^N$  as a 2-tuple  $\langle P, L \rangle$  consisting of a floating point scalar  $P$  and an integer level indicator  $L$ .  $P$  represents the “goodness” of  $\mathbf{x}$  at level  $L$ .



## Computing Performance

- To determine  $\langle P, L \rangle$ 
  - sum the constraint violations in each level
  - $L$  is assigned to the first level, beginning at level 1, to have any violation and  $P$  is assigned the sum of the violations at  $L$
  - if no violations occur,  $L \leftarrow W$  and  $P \leftarrow F(\mathbf{x})$



## Comparing Performances

- The partitioned performances of two locations  $\mathbf{x}_1$  ( $\langle P_1, L_1 \rangle$ ) and  $\mathbf{x}_2$  ( $\langle P_2, L_2 \rangle$ ) are compared as follows:
  - if ( $L_1 == L_2$ )
    - if ( $P_1 > P_2$ )  $\mathbf{x}_1$  is better, otherwise  $\mathbf{x}_2$  is better
  - if ( $L_1 > L_2$ )
    - $\mathbf{x}_1$  is better
  - if ( $L_2 > L_1$ )
    - $\mathbf{x}_2$  is better



## SPIDER Method

- Applies partitioned performances to a new variation of the Nelder and Mead polytope algorithm
- Rules for ordering and movement using partitioned performances are demonstrated

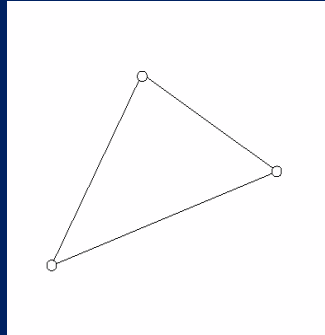


## What is a “SPIDER”?

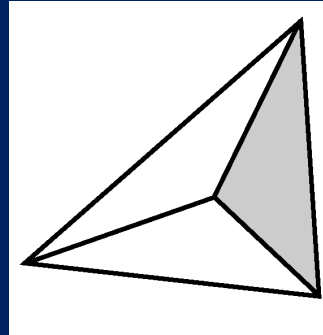
- Assuming we are maximizing an  $n$ -dimensional objective function  $F$ , SPIDER consists of  $n+1$  “legs”, where
  - each leg contains its position in space
  - associated with each leg is a partitioned performance



## What is a “SPIDER”?



When  $n = 2$ , a triangle

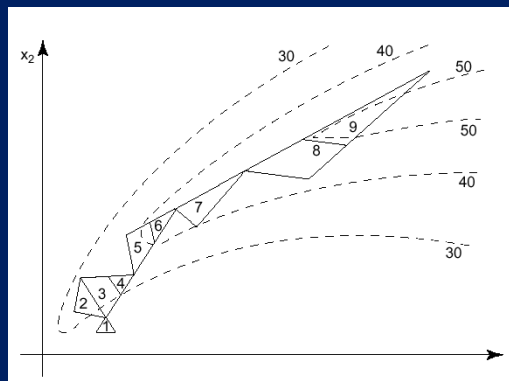


When  $n = 3$ , a tetrahedron



## What does SPIDER do?

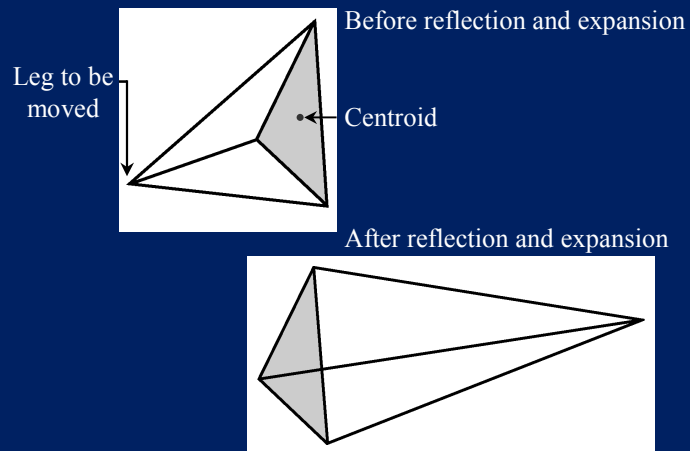
- Crawl: each leg is at a known “elevation” on the performance “hill”, and it is SPIDER’s task to crawl up the hill and maximize performance





## How SPIDER walks

- By moving each leg through the centroid of the remaining legs



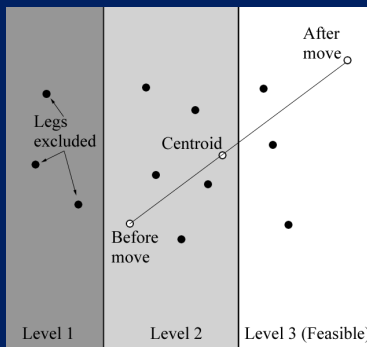
## How SPIDER walks

- Repeat N times
  - Sort legs of SPIDER, from worst to best. Label worst and best legs.
  - For each leg L, in worst to best order
    - Determine centroid
    - Compute position and performance of a trial leg,  $L_{\text{trial}}$ 
      - if L is not the best leg, reflect and expand *through* centroid
      - if L is the best leg, reflect and expand *away* from centroid
    - If move successful, accept trial, relabel worst and best leg if required
  - EndFor
  - Shrink SPIDER if best leg has not improved
  - Rebuild SPIDER if successive shrinks exceed threshold
- EndRepeat



## Rules for centroid computation

- Exclude leg being moved (L)
- Exclude legs at a lower level than L
  - ☞ this helps to give SPIDER a better sense of direction along constraint boundaries



## Rules for moving a non-best leg

- Same level (level of  $L_{\text{trial}}$  == level of L)
  - ☞ accept trial leg if
    - P value of  $L_{\text{trial}}$  > P value of L
- Going down levels (level of  $L_{\text{trial}}$  < level of L)
  - ☞ accept trial leg if its better than the worst leg
- Going up levels (level of  $L_{\text{trial}}$  > level of L)
  - ☞ accept trial leg if its better than the best leg



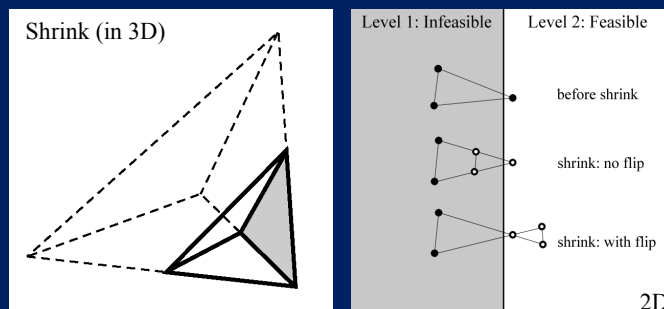
## Rules for moving the best leg

- It must improve in performance in order to move
- This gives SPIDER the ability to “straddle” and thus track along a constraint boundary



## Rules for shrinking SPIDER

- Shrink the vertices at the same level as the best leg toward the best leg, and flip (as well as shrink) vertices at lower levels over the best leg
- Flipping helps to move legs across a constraint boundary towards feasibility

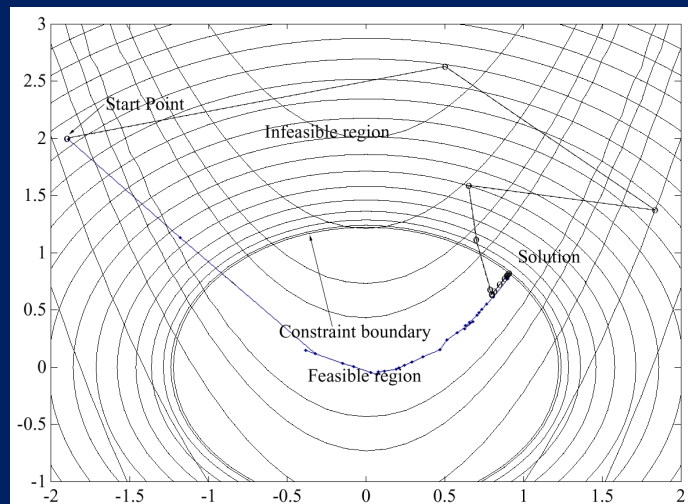


## A Matlab Test Problem

- Sequential Quadratic Programming (SQP) methods represent the state-of-the-art in nonlinear constrained optimization
- SQP methods out perform every other tested method in terms of efficiency, accuracy, and percentage of successful solutions, over a large number of test problems
- On a Matlab test problem
  - Matlab SQP Implementation, 96 function calls
  - SPIDER, 108 function calls



## A Matlab Test Problem



SPIDER walk in blue, SQP walk in black



[illegible]